

Lecture 8

CS 6260 Automated Planning and Learning

Heuristics for Planning 2

Department of Computer Science and Engineering

Indian Institute of Technology Madras



Recap: Heuristics leveraging PDDL

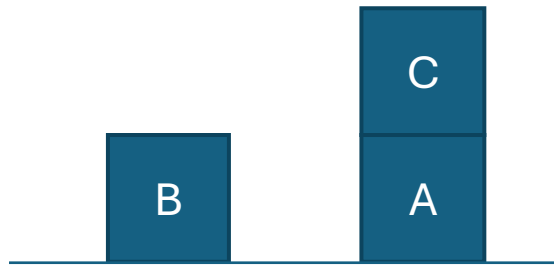
Ideas for Generating Heuristics and Search Control

- Goal decomposition (not quite)
- Planning graph
- Landmarks

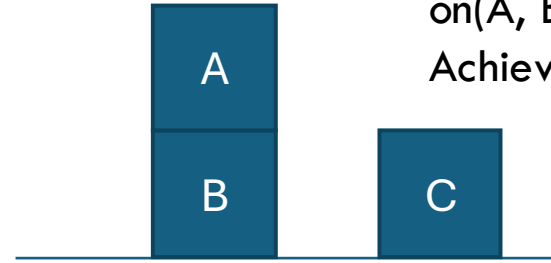
Goal Decomposition: Sussman's Anomaly

Goal requires: $\text{on}(A, B)$; $\text{on}(B, C)$

What happens if we try to achieve them independently?

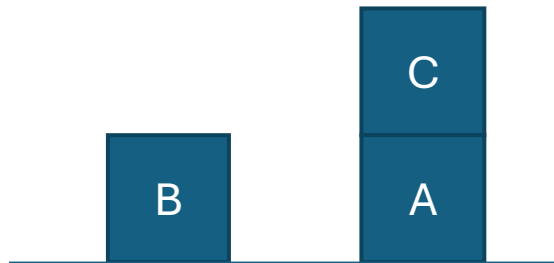


Try to achieve: $\text{on}(A, B)$
 $\text{Move}(C, \text{table}); \text{Move}(A, B)$

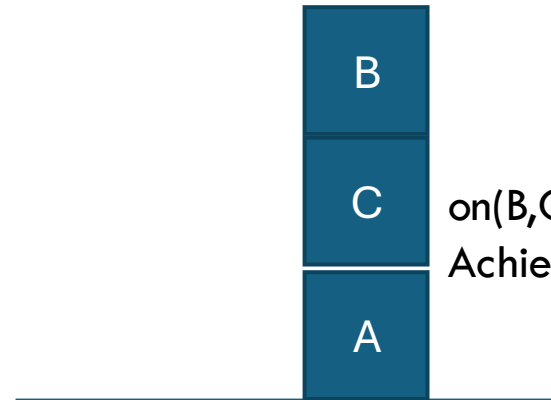


$\text{on}(A, B)$ achieved 😊

Achieving $\text{on}(B, C)$ will undo it ☹️



Try to achieve: $\text{on}(B, C)$
 $\text{Move}(B, C)$



$\text{on}(B, C)$ achieved 😊

Achieving $\text{on}(A, B)$ will undo it ☹️

Planning Graph

$Init(Have(Cake))$

$Goal(Have(Cake) \wedge Eaten(Cake))$

$Action(Eat(Cake))$

PRECOND: $Have(Cake)$

EFFECT: $\neg Have(Cake) \wedge Eaten(Cake)$

$Action(Bake(Cake))$

PRECOND: $\neg Have(Cake)$

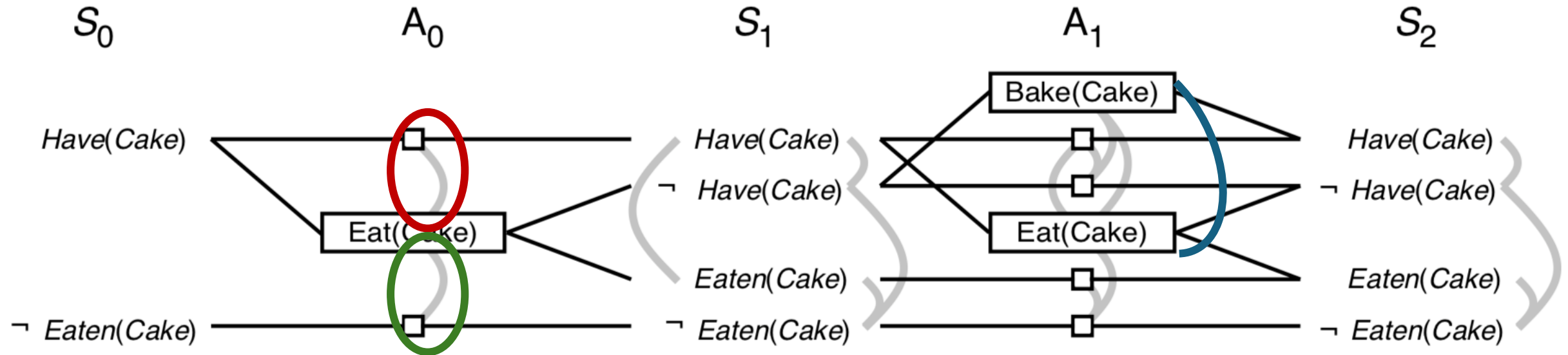
EFFECT: $Have(Cake)$

Actions in a layer are *mutex* if:

Inconsistent Effects: one action's effect negates another's effect

Interference: one action negates another's precondition

Competing needs: one action requires negation of another's precondition



Planning Graph

$Init(Have(Cake))$

$Goal(Have(Cake) \wedge Eaten(Cake))$

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PRECOND: $Have(Cake)$

EFFECT: $\neg Have(Cake) \wedge Eaten(Cake)$

$Action(Bake(Cake))$

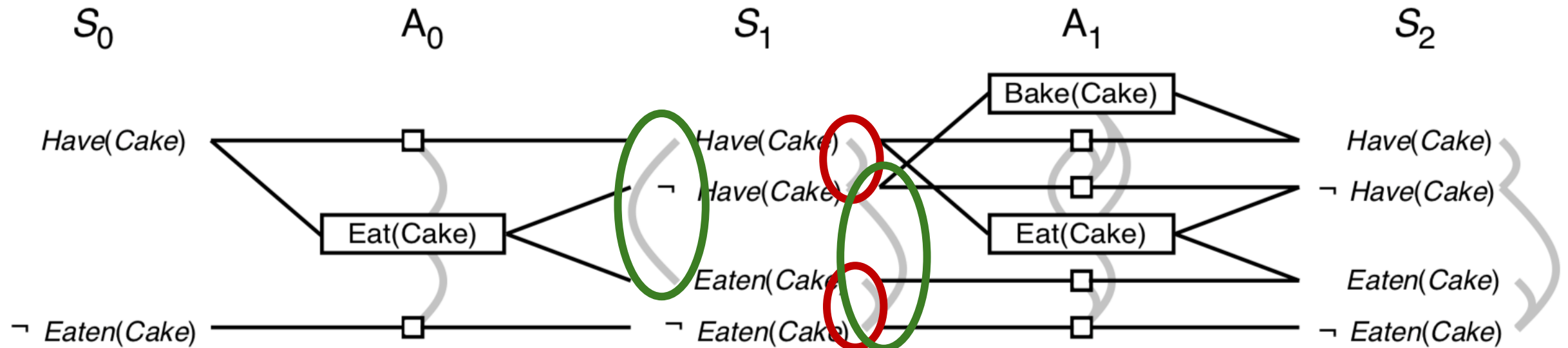
PRECOND: $\neg Have(Cake)$

EFFECT: $Have(Cake)$

Propositions in a layer are *mutex* if:

Negation: one negates the other

Inconsistent Support: all ways of achieving them are pairwise mutex



function GRAPHPLAN(*problem*) **returns** solution or failure

graph $\leftarrow$ INITIAL-PLANNING-GRAPH(*problem*)

goals $\leftarrow$ GOALS[*problem*]

loop do

if *goals* all non-mutex in last level of *graph* **then do**

solution $\leftarrow$ EXTRACT-SOLUTION(*graph*, *goals*, LENGTH(*graph*))

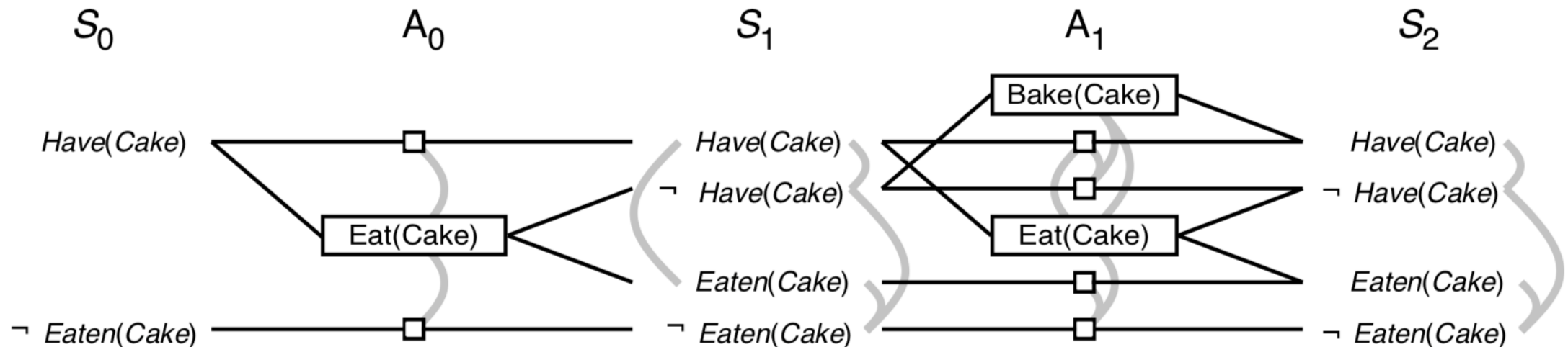
if *solution* $\neq$ failure **then return** *solution*

else if NO-SOLUTION-POSSIBLE(*graph*) **then return** failure

graph $\leftarrow$ EXPAND-GRAPH(*graph*, *problem*)

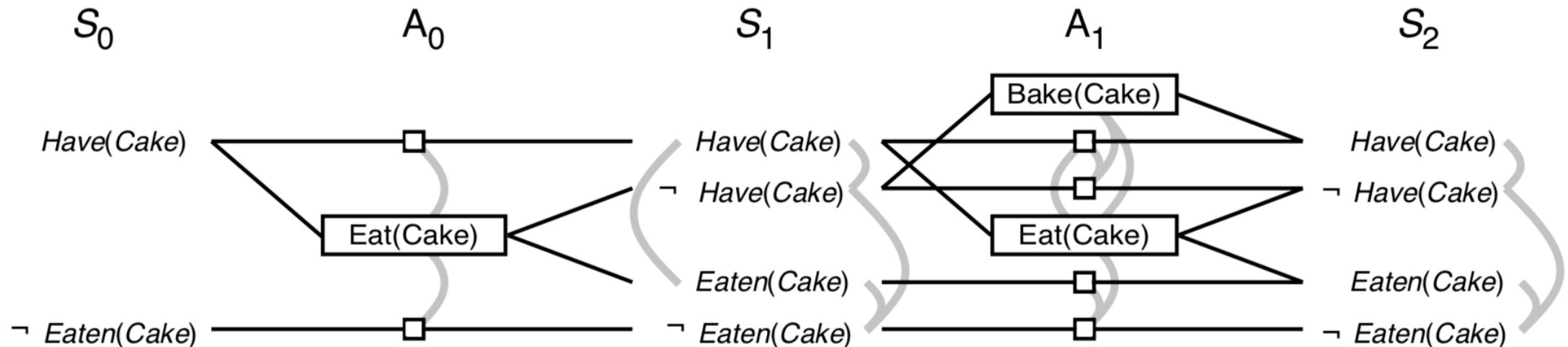
Extracting Heuristics From the Planning Graph

- Mutex relations can be used to try to extract plans that achieve the goal
- It is possible that a plan is not feasible within a given expansion of graphplan
 - Add another layer, try again
 - Until the graph levels off
 - If all the goal literals do not appear in the final level, goal is unreachable



Extracting Heuristics From the Planning Graph

- Expand the graph until you reach a fixed point
- If a literal doesn't occur in the last layer, it cannot be achieved
- Level-cost heuristic
 - Estimated cost of achieving a goal literal = **level at which it first appears**
 - Is this admissible?
- What about goals that include multiple literals?
 - Max-level heuristic
 - Level-sum heuristic
 - Is this optimistic (admissible) or pessimistic (non-admissible)?



Heuristic Search Planner (HSP)

- Consider only positive effects of actions (ignore negated literals in effects)
 - This creates a relaxed problem: it allows more actions to be applicable in a state than truly possible
 - Recall: To generate a heuristic, use cost of solution in relaxed problem...
 - But computing optimal solutions is too hard even for such relaxed problems
- Compute estimated cost of achieving an atom as

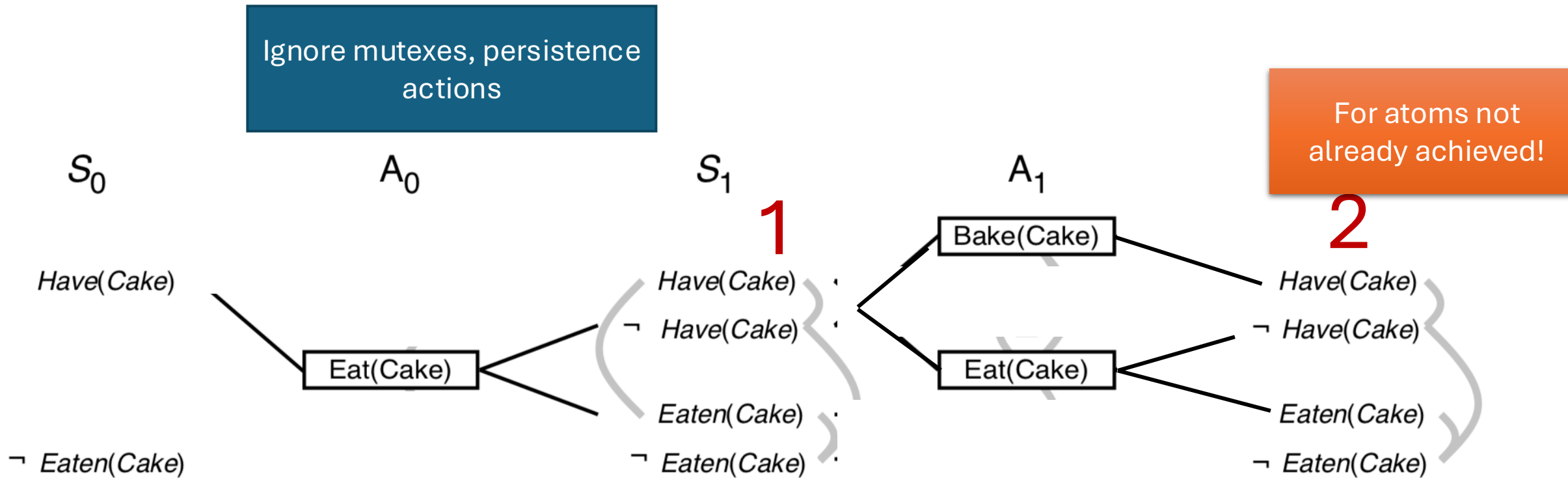
$$g_s(p) = \begin{cases} 0 & \text{if } p \in s, \\ \min_{op \in O(p)} [1 + g_s(Prec(op))] & \text{otherwise,} \end{cases}$$

Operators that add p

Estimated cost of achieving preconditions of op

Heuristic Search Planner (HSP)

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Heuristic Search Planner (HSP)

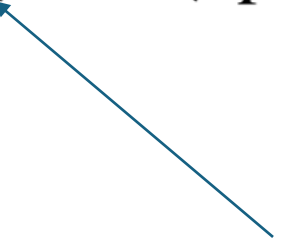
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Options for g_s of a set of atoms:

$$g_s^{min}(C) = \min_{r \in C} g_s(r) \qquad g_s^{max}(C) = \max_{r \in C} g_s(r) \qquad g_s^+(C) = \sum_{r \in C} g_s(r)$$

Heuristic Search Planner (HSP)

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Estimated cost of achieving preconditions of op

Use this to compute the heuristic estimate:

$$h(s) := g_s(G)$$

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Use this to compute the heuristic estimate:

$$h(s) := g_s(G)$$

h_{min} uses g_s^{min}

h_{max} uses g_s^{max}

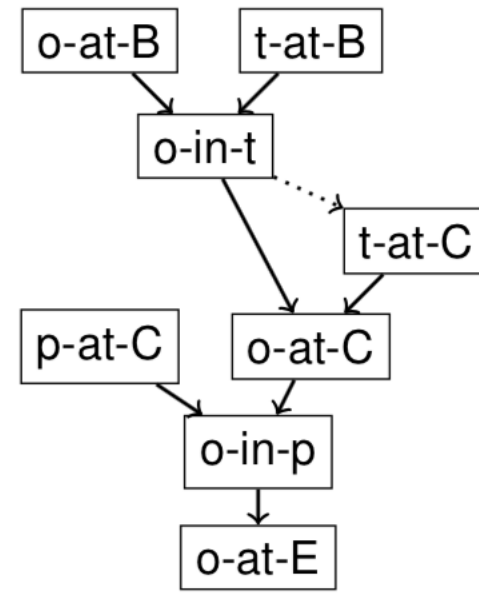
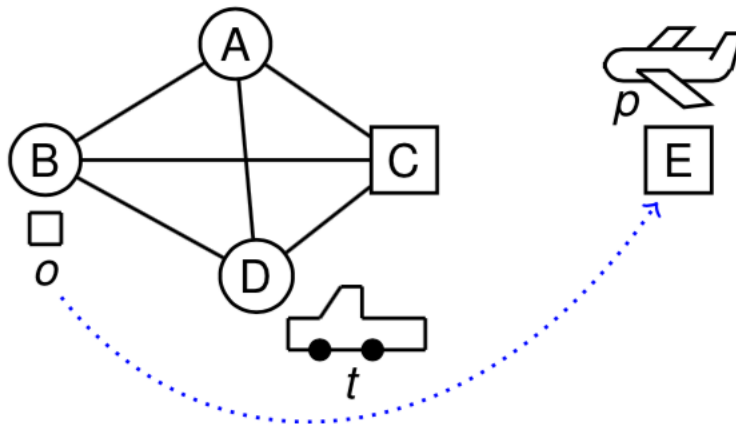
h_{add} uses g_s^+

FF Planner

- <https://arxiv.org/abs/1106.0675>
- <https://fai.cs.uni-saarland.de/hoffmann/ff.html>
- Uses custom h_{FF} heuristic.
- h_{FF} is not admissible but practically better than h_{max} .
- $h_{min} \leq h_{max} \leq h_{FF} \leq h_{add}$.

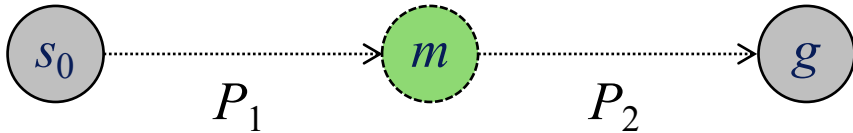
Landmarks in Planning

- A *landmark* is a formula that must become true at some point in every plan
- An *action landmark* is an action that must occur in every plan
- Landmarks and action landmarks can be ordered



Why are Landmarks Useful?

- Can break a problem down into smaller subproblems



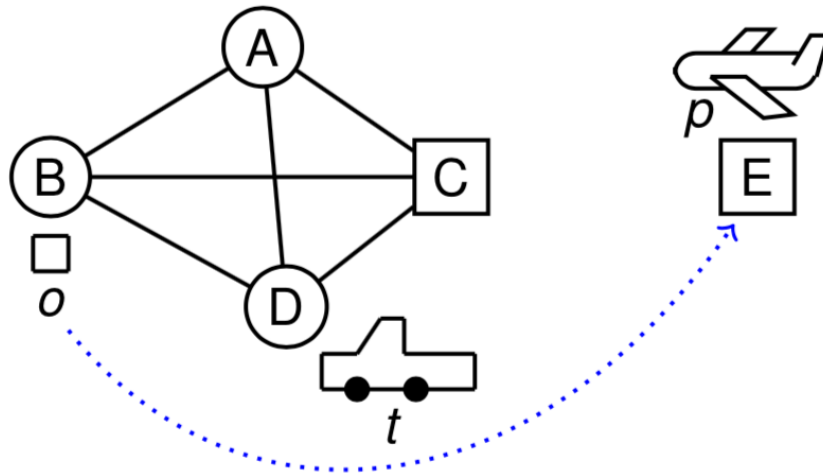
- Suppose m is a landmark
 - Every solution to P must achieve m
- Possible strategy:
 - find a plan to go from s_0 to any state s_1 that satisfies m
 - find a plan to go from s_1 to any state s_2 that satisfies g



- Suppose m_1, m_2, m_3 are landmarks
 - Every solution to P must achieve m_1 , then m_2 , then m_3
- Possible strategy:
 - find a plan to go from s_0 to any state s_1 that satisfies m_1
 - find a plan to go from s_1 to any state s_2 that satisfies m_2
 - ...

Computing Landmarks

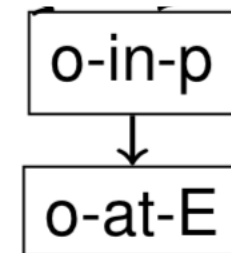
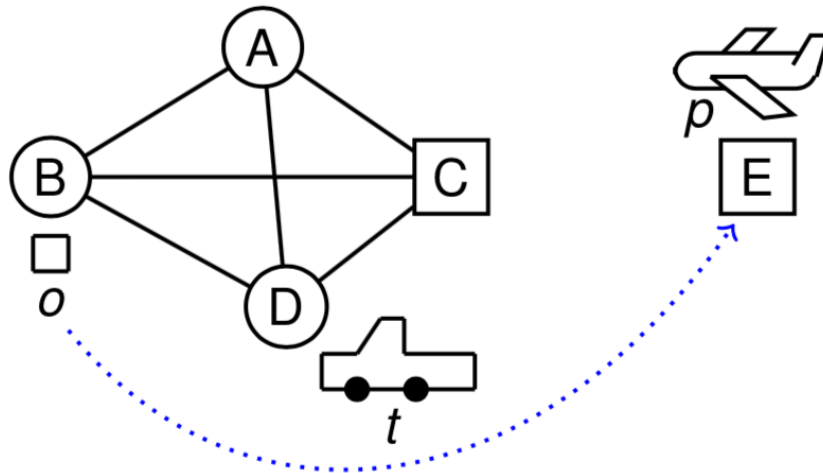
- Backchaining:
 - Every goal is a landmark
 - If B is a landmark and A is a common precondition in all actions that achieve B ,
 - A is a landmark
 - $A \preceq B$



o-at-E

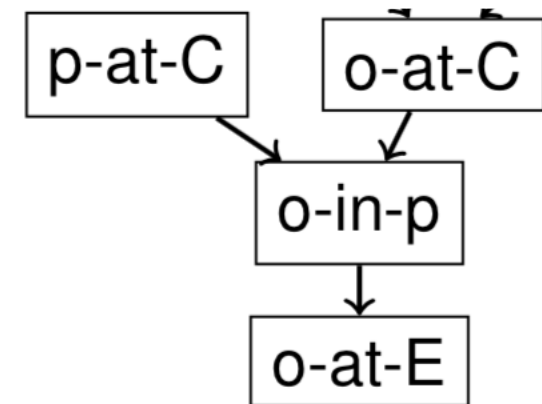
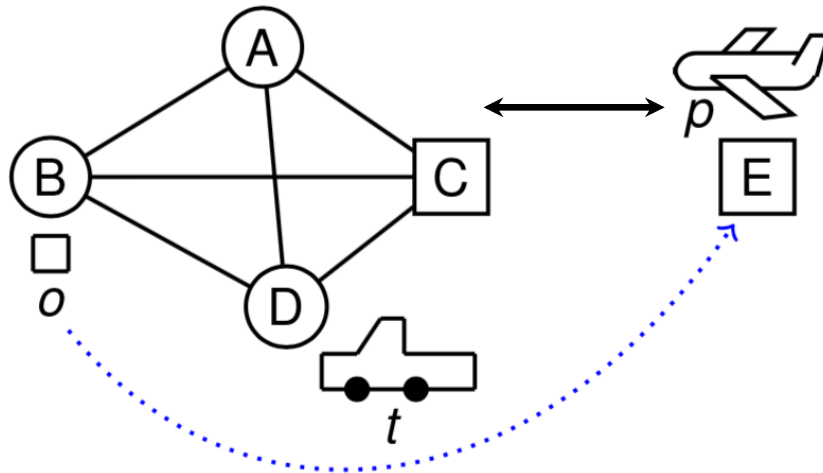
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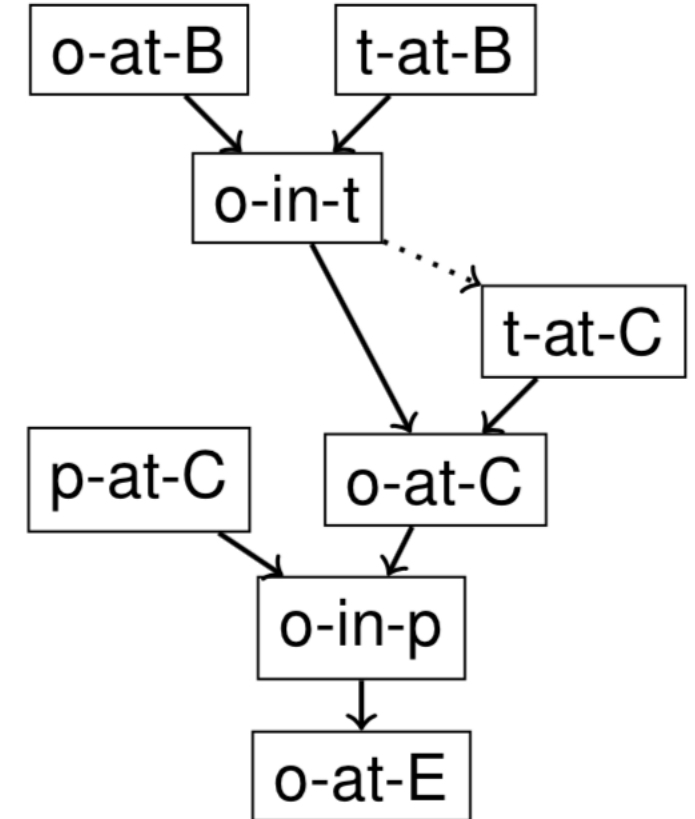
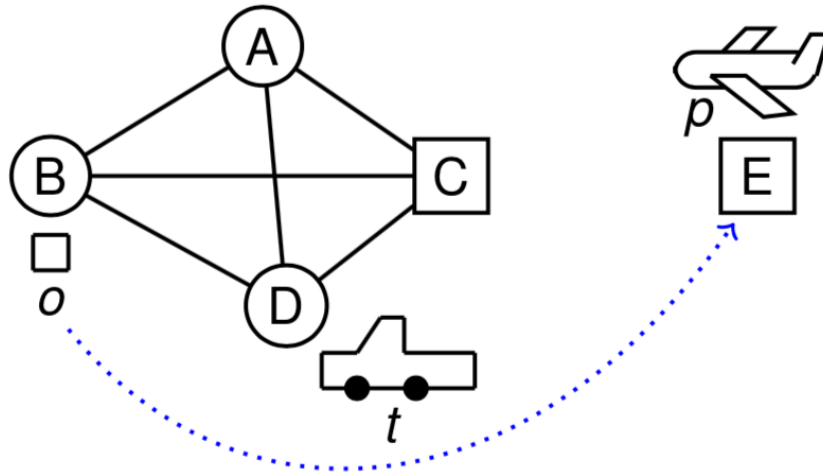
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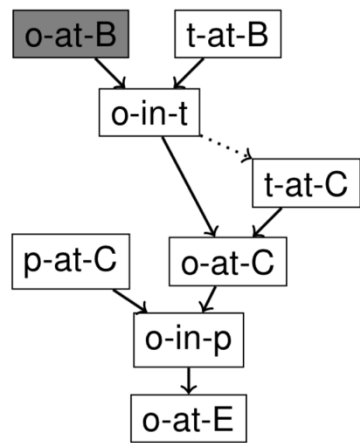
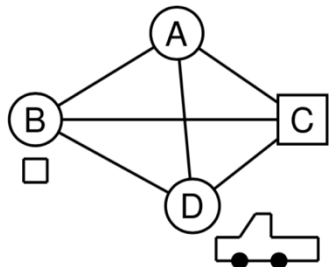
Computing Landmarks

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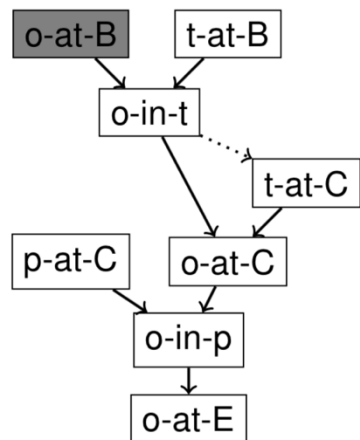
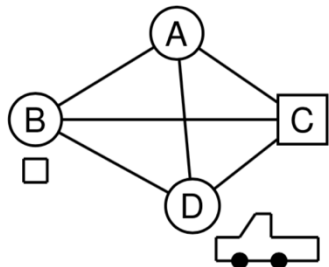


Using Landmarks as Heuristics

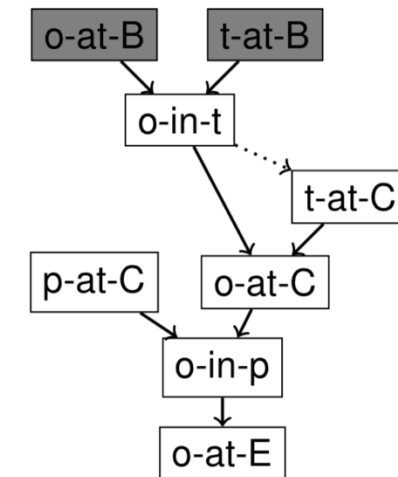
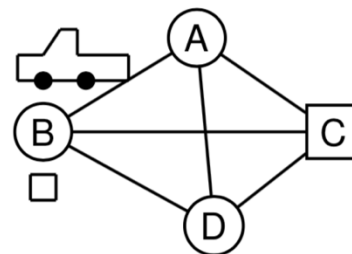
- Two ways of using landmarks: subgoals, components in heuristic estimate
- Subgoals:
 - Select earliest landmark(s)
 - Set as disjunctive goal for any planner
 - After plan is found, repeat



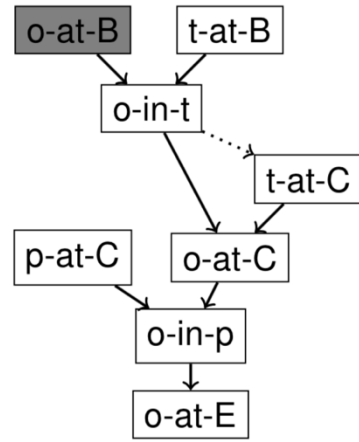
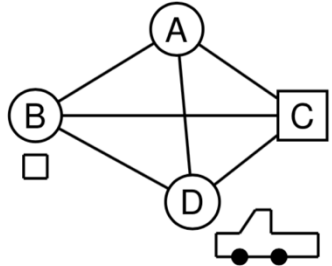
- Partial plan: $\emptyset$
- Goal: $t\text{-at-B} \vee p\text{-at-C}$



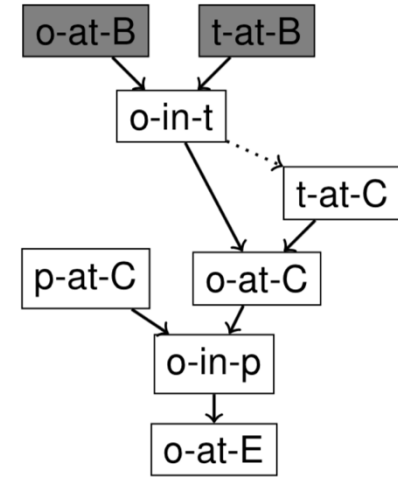
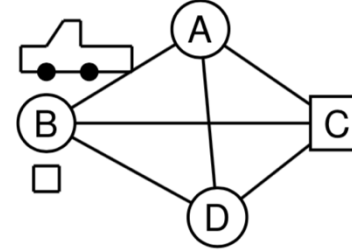
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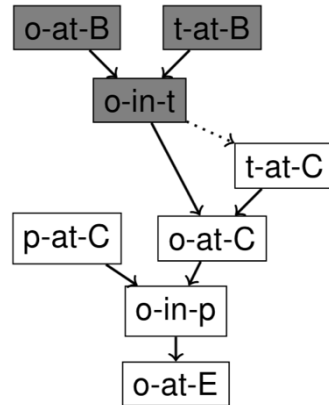
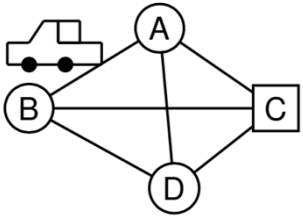
- Partial plan: Drive-t-B
- Goal: $o\text{-in-t} \vee p\text{-at-C}$



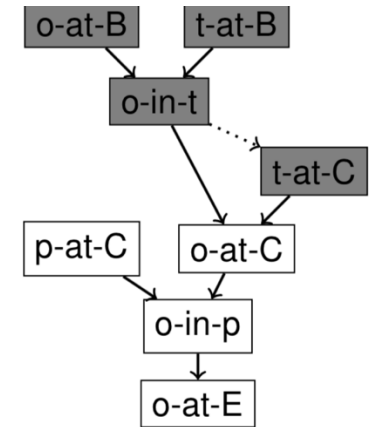
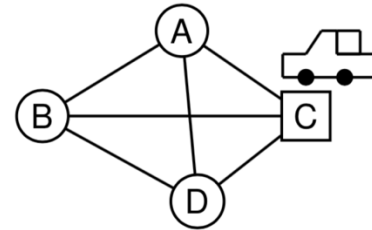
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- Goal: $t\text{-at-B} \vee p\text{-at-C}$



- Partial plan: Drive-t-B
- Goal: $o\text{-in-t} \vee p\text{-at-C}$



- Partial plan: Drive-t-B, Load-o-B
- Goal: $t\text{-at-C} \vee p\text{-at-C}$



- Partial plan: Drive-t-B, Load-o-B, Drive-t-C
- Goal: $o\text{-at-C} \vee p\text{-at-C}$

Using Landmarks

- Using landmarks as subgoals is not guaranteed to be efficient
 - Order of achieving mutually unordered landmarks can affect size of the plan
 - Variant of sussman's anomaly
 - May fail for problems with deadends
- Alternative:
 - Use # unachieved landmarks as a part of the heuristic estimate

Backward Search

- Forward search: forward from initial state
 - In state s , choose applicable action a
 - Compute state transition $s' = T(s, a)$
- Backward search: backward from the goal
 - For goal g , choose *relevant* action a
 - A possible “last action” before the goal
 - Sometimes this has a lower branching factor
- Compute *inverse* state transition $g' = T^{-1}(g, a)$
 - $g' =$ properties a state s' should satisfy in order for $T(s', a)$ to satisfy g
- Equivalently, if $S_g = \{\text{all states that satisfy } g\}$ then
 - $S_{g'} = \{\text{all states } s \text{ such that } T(s, a) \in S_g\}$